

Constructions for 4-Polytopes and the Cone of Flag Vectors

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ABSTRACT. We describe a construction for d -polytopes generalising the well known stacking operation. The construction is applied to produce 2-simplicial and 2-simple 4-polytopes with $g_2 = 0$ on any number of $n \geq 13$ vertices. In particular, this implies that the ray ℓ_1 , described by Bayer (1987), is fully contained in the convex hull of all flag vectors of 4-polytopes. Especially interesting examples on 9, 10 and 11 vertices are presented.

1. Introduction

It is a wide open problem in discrete geometry to understand the combinatorial properties of polytopes, which can be described as the convex hull of finitely many points in some \mathbb{R}^n . Even the more special task to characterise the f - and flag vectors of polytopes (and more generally of spheres) of arbitrary dimension seems to be very hard. For 3-dimensional polytopes, Steinitz [Ste06] gave a complete characterisation. For arbitrary dimension, the problem is still open, although some conditions are known. In particular, all linear relations between the entries of flag vectors are described by the Generalized Dehn-Sommerville equations [BB85] and the admissible f -vectors of simplicial polytopes are classified in terms of their g -vectors by the g -theorem of Billera, Lee and Stanley and McMullen (see for instance [Zie95, Thm. 8.35]). Additionally, for dimension 4 a linear approximation of the set of flag vectors and f -vectors was given by Bayer [Bay87]. A different view on the f -vector characterisation was provided by Ziegler [Zie02]. Since then, some progress has been made (cf. [PZ04], [Zie04]), but it is still not known what the linear cone of f - resp. flag vectors looks like.

In this paper we introduce a new polytope construction method, examine some of its combinatorial properties and construct 4-polytopes with arbitrarily high numbers of vertices that are extremal for the flag vector cone. In particular, we prove the following two theorems (see below for definitions).

2000 *Mathematics Subject Classification.* Primary 52B05;52B12.

Key words and phrases. Polytope, Flag Vector, Flag Vector Cone, g -Vector, Elementary Polytope, 2-Simple, 2-Simplicial.

The first author was supported by the Deutsche Forschungsgemeinschaft within the European graduate program ‘Combinatorics, Geometry, and Computation’ (No. GRK 588/2).

The second author was supported by the DFG Leibniz grant of G. M. Ziegler.

THEOREM 4.3. *Elementary 2-simplicial, 2-simple 4-polytopes with k vertices exist for $k = 5, 9, 10, 11$ and $k \geq 13$.*

This implies (using the notation of Bayer [Bay87]):

COROLLARY 4.4. *The ray ℓ_1 is contained in the convex hull of all flag vectors of 4-polytopes.*

Additionally, we briefly analyse the consequences of the various recent polytope constructions in [PZ04] and [Zie04] for the flag vector cone.

Acknowledgements. The authors would like to thank Günter M. Ziegler, who supported this work at several occasions and in particular pointed out the significance of our series of polytopes, Eran Nevo and Raman Sanyal for helpful discussions, and the referee for several suggestions improving the exposition.

1.1. General preliminaries. We first give the basic definitions as well as an overview over important related results.

Let P be a d -polytope and $[d] = \{0, \dots, d-1\}$. The *f-vector* of P is the d -dimensional vector

$$f(P) = (f_0, f_1, \dots, f_{d-1}),$$

where f_i for $0 \leq i \leq d-1$ denotes the number of i -dimensional faces of P . The *flag vector* of P is the 2^d -dimensional vector

$$\text{flag}(P) = (f_S)_{S \subseteq [d]},$$

where f_S for a subset $S = \{i_1, \dots, i_k\}$ of $[d]$ denotes the number of face chains $\emptyset \subset F_{i_1} \subset \dots \subset F_{i_k} \subset P$ such that $\dim F_{i_j} = i_j$ for $1 \leq j \leq k$. We usually write $f_{i_1 i_2 \dots i_k}$ instead of $f_{\{i_1, i_2, \dots, i_k\}}$. Faces of codimension 1, 2 and 3 are called *facets*, *ridges* and *subridges* respectively.

The Generalized Dehn-Sommerville equations by Bayer and Billera [BB85] imply that the flag vectors of d -polytopes lie in an $(F_d - 1)$ -dimensional affine subspace of \mathbb{R}^{2^d} , where $(F_k)_{k \geq 0} = (1, 1, 2, 3, \dots)$ is the series of Fibonacci numbers.

P is *k -simplicial* ($0 \leq k \leq d-1$) if all its k -dimensional faces are k -simplices, i.e. contain exactly $k+1$ vertices; equivalently, if in the face lattice of P all intervals $[\emptyset, F]$ for $\dim F = k$ are boolean.

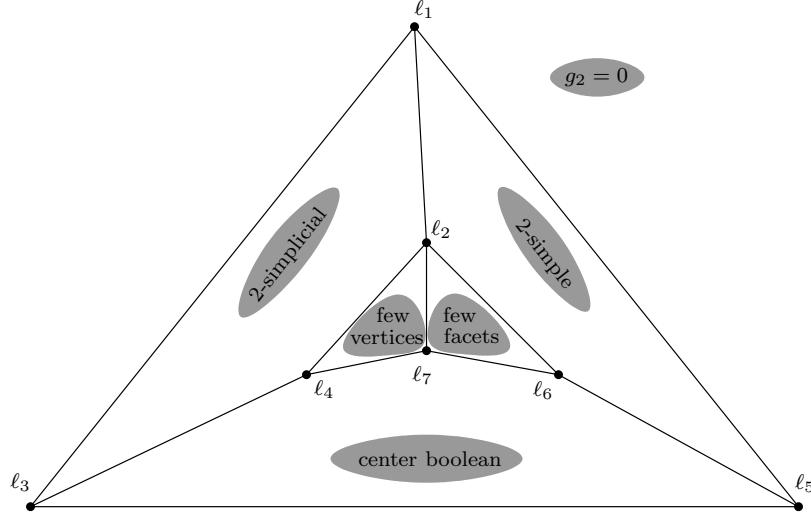
P is *h -simple* ($0 \leq h \leq d-1$) if the dual polytope P^Δ is h -simplicial; equivalently, if all intervals $[F, P]$ with $\dim F = d-h-1$ are boolean or equivalently, if all its $(d-h-1)$ -dimensional faces are contained in exactly $h+1$ facets.

The *g-vector* of a polytope can be defined in general using generating functions (see [Sta97]); its entries can be written as linear combinations of the entries of the flag vector. We will, however, focus on one special entry.

DEFINITION 1.1. For a d -polytope P define

$$(1) \quad g_2(P) = f_{02} - 3f_2 + f_1 - df_0 + \binom{d+1}{2}.$$

It has been shown by Kalai [Kal87] via rigidity theory that $g_2(P) \geq 0$ for every d -polytope P with $d \geq 4$. Polytopes P with $g_2(P) = 0$ are called *elementary*. It is an interesting open problem to characterise all elementary polytopes. See also [Kal94] for a survey on this topic.

FIGURE 1. Hyperplane section through the cone \mathcal{D} , according to [Bay87]

1.2. Flag vectors of 4-polytopes. In the last two sections we will focus on 4-polytopes, where the situation is understood slightly better than in general. It is easy to prove that for 2-simplicial, 2-simple 4-polytopes the f -vector is symmetric and the flag vector is completely determined by the f -vector.

Bayer [Bay87] described all the known linear inequalities for flag vectors of 4-polytopes. Let $\mathcal{D} \subseteq \mathbb{R}^4$ be the polyhedron defined by these inequalities and $\mathcal{C} \subseteq \mathcal{D}$ the convex hull of all flag vectors of 4-polytopes. Then \mathcal{D} is a 4-dimensional cone with the flag vector of the 4-simplex as its apex. Figure 1 illustrates a hyperplane section through this cone. It can be viewed as a 3-polytope with its vertices, edges and faces representing special properties of 4-polytopes, in the sense that polytopes whose flag vectors lie on the respective faces have these properties. The main question is how close the cone \mathcal{D} approximates \mathcal{C} . The task is therefore to find examples of polytopes with flag vectors in extremal regions of \mathcal{D} .

In Table 1 we give a summary of what is known for the rays of \mathcal{D} ; note that ℓ_4 and ℓ_6 are not contained in \mathcal{C} itself, but in its closure (cf. [Bay87, Sec. 2]).

Until now only two flag vectors of polytopes on ℓ_1 were known, that of the 4-simplex Δ_4 and that of the 4-dimensional hypersimplex $\Delta_4(2)$ (for a definition

TABLE 1. Known polytopes on or close to the rays of \mathcal{D}

ray	property	examples
ℓ_1	2-simplicial, 2-simple 4-polytopes with $g_2 = 0$	see Theorem 4.3
ℓ_2	‘fat’ 2-simplicial, 2-simple 4-polytopes	unknown
ℓ_7	‘fat’ center boolean 4-polytopes	unknown
ℓ_4	simplicial 4-polytopes with few vertices	cyclic polytopes
ℓ_3	simplicial 4-polytopes with $g_2 = 0$	stacked polytopes
ℓ_5	simple 4-polytopes with $g_2 = 0$	truncated polytopes
ℓ_6	simple 4-polytopes with few facets	dual cyclic polytopes

see [Zie95, Ch. 0] for instance). Hence it was clear that ℓ_1 contained an edge of \mathcal{C} , but not, whether ℓ_1 is a ray of \mathcal{C} , nor if ℓ_1 contained any further flag vectors at all. We establish in Section 4 that ℓ_1 is indeed an extremal ray of \mathcal{C} .

A few more regions of \mathcal{D} deserve a closer study. The set of flag vectors of general 2-simplicial, 2-simple 4-polytopes is a subset of the 2-dimensional cone spanned by ℓ_1 and ℓ_2 . It is indeed a 2-dimensional set – there are 2-simplicial, 2-simple 4-polytopes with the same number of vertices, but different numbers of edges (cf. [PZ04]). Analogously, the set of flag vectors of elementary 4-polytopes is contained in the 3-dimensional cone spanned by ℓ_1 , ℓ_3 and ℓ_5 . There are also a number of extreme examples: stacked polytopes, iterated pyramids over n -gons and multiplexes (cf. [BBS02]). Also, the constructions described in the next sections produce in general elementary polytopes when applied to such, as can be seen from Corollary 3.8. Note that 2-simplicity, 2-simpliciality and $g_2 = 0$ are three independent properties, that is, there are 4-polytopes with any combination of these properties.

Bayer conjectured that the hyperplane determined by ℓ_2 , ℓ_4 and ℓ_6 yields a valid inequality for 4-polytopes. This is not true, since there exist polytopes with flag vectors close to the interior of the edges $[\ell_4, \ell_7]$, and dually $[\ell_6, \ell_7]$, as shown by Joswig & Ziegler [JZ00] and Ziegler [Zie04]. However, a hyperplane cutting off ℓ_7 may still be possible. Extremal polytopes known in this respect are projected products of polygons by Ziegler [Zie04] and polytopes obtained by the E -construction (Paffenholz & Ziegler [PZ04]). The most restrictive linear inequality that is compatible with these examples would then be

$$f_{03} - 140 \geq 4(f_1 + f_2) - 20(f_0 + f_3).$$

In terms of the *fatness* F and *complexity* C of 4-polytopes, as introduced by Ziegler [Zie02], this inequality reads $4F - C \leq 20$. In this respect, providing an upper bound for fatness, that is, bounding the number of edges and ridges by the number of vertices and facets, would be helpful.

2. General Construction

This section consists of two parts. In the first part we provide the basic tool for the construction of a family of 2-simple and 2-simplicial 4-polytopes with vanishing g_2 . The operation has many more applications than the ones we will discuss in more detail in the second part. Hence, we give all definitions and theorems for arbitrary dimension $d \geq 3$. We indicate some of the additional applications we have encountered so far.

In the second part we examine settings in which the operation can be applied, determine the facet types that can occur in the construction and prove that some interesting properties of polytopes are preserved. This part will be much more specifically tailored for what we need for the polytope families defined later, as some of the f -vector computations tend to get complicated in a more general setting.

In Section 3 we combine several instances of our basic tool to obtain two special constructions \mathcal{I}^1 and \mathcal{I}^2 producing 4-polytopes $\mathcal{I}^1(P; S)$ and $\mathcal{I}^2(P; S)$ out of a polytope P and a facet S of P . Using the properties of the construction discussed in this section we prove that 2-simplicity, 2-simpliciality and the value of g_2 is preserved by the construction.

2.1. Pseudo-Stacking. The operation we introduce here is a generalisation of the well known stacking operation. In most cases it adds one new vertex to the polytope. We need some notation for this.

DEFINITION 2.1. Let P be a d -polytope. A *simplex facet* of P is a facet of P that is combinatorially equivalent to a $(d - 1)$ -simplex.

Let $H := \{x \in \mathbb{R}^d \mid \langle x, v \rangle = \ell\}$ for some $v \in \mathbb{R}^d$ and $\ell \in \mathbb{R}$ be an affine hyperplane. By $H^+ := \{x \in \mathbb{R}^d \mid \langle x, v \rangle > \ell\}$ we denote the positive half space defined by H and similarly the negative half space by H^- .

Let F be a facet of a d -polytope P . We denote by H_F the unique affine hyperplane that contains F , oriented in such a way that P is contained in $H_F^+ \cup H_F$. By $\text{Fac}(P)$ we denote the set of all facets of P , by $\mathcal{H}(P) := \{H_F \mid F \in \text{Fac}(P)\}$ the set of all hyperplanes coming from the facets, and by $\mathcal{H}^+(P)$ the set $\{H_F^+ \mid F \in \text{Fac}(P)\}$ of half-spaces.

Further let $\text{adj}(F)$ be the set of facets of P adjacent to F . For a subset \mathcal{F} of $\text{adj}(F)$ we denote the set of hyperplanes defined by the facets in \mathcal{F} by $\mathcal{H}_{\mathcal{F}}$, and the set of positive half-spaces determined by the hyperplanes in $\mathcal{H}_{\mathcal{F}}$ by $\mathcal{H}_{\mathcal{F}}^+$.

Finally, if v is some point in \mathbb{R}^d then v lies *beyond* a facet F if $v \in H_F^-$ and v lies *beneath* F if $v \in H_F^+$.

Let S be a facet of a d -polytope P and \mathcal{F}, \mathcal{N} be disjoint subsets of $\text{adj}(S)$. We define the region $\mathcal{R}_{\mathcal{F}, \mathcal{N}}^S(P)$ in \mathbb{R}^d by

$$\mathcal{R}_{\mathcal{F}, \mathcal{N}}^S(P) := \left(\bigcap_{H \in \mathcal{A}} H^+ \right) \cap \left(\bigcap_{H \in \mathcal{H}_{\mathcal{F}}} H \right) \cap \left(\bigcap_{H \in \mathcal{H}_{\mathcal{N}} \cup \{H_S\}} H^- \right)$$

for $\mathcal{A} := \mathcal{H}(P) \setminus (\mathcal{H}_{\mathcal{F}} \cup \mathcal{H}_{\mathcal{N}} \cup \{H_S\})$.

DEFINITION 2.2 (Pseudo-Stacking). Let P be a d -polytope and S a simplex facet of P . Choose two disjoint sets $\mathcal{F}, \mathcal{N} \subseteq \text{adj}(S)$. Assume that $\mathcal{R}_{\mathcal{F}, \mathcal{N}}^S(P) \neq \emptyset$.

The *pseudo-stacking* $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ of P above S with respect to \mathcal{F} and \mathcal{N} is the convex hull $\text{conv}(P \cup v)$ of P with a point $v \in \mathcal{R}_{\mathcal{F}, \mathcal{N}}^S(P)$. If $\mathcal{N} = \emptyset$ then we omit it in the notation and write $\mathbf{PS}_{\mathcal{F}}^S(P)$ for the pseudo-stacking of P above S .

Compare this definition to Grünbaum [Grü03, Section 5.2]. Also, a similar concept was studied by Altshuler and Shemer [AS84] for the purpose of enumeration of 4-polytopes with few vertices.

In plain words, a point $v \in \mathcal{R}_{\mathcal{F}, \mathcal{N}}^S(P)$ lies *beyond* S and all facet hyperplanes coming from facets in \mathcal{N} , it lies *in* all facet hyperplanes from facets in \mathcal{F} , and *beneath* all other facet hyperplanes. A priori, the set $\mathcal{R}_{\mathcal{F}, \mathcal{N}}^S(P)$ need not contain a point at all, but we will show some conditions that guarantee that this set is non-empty. See Figures 2–5 for some illustrations of this operation with various choices of \mathcal{F} and \mathcal{N} . We have the following simple fact about the pseudo-stacking operation:

PROPOSITION 2.3. *The combinatorial properties of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ do not depend on the actual choice of the point $v \in \mathcal{R}_{\mathcal{F}, \mathcal{N}}^S(P)$.* \square

REMARK 2.4. The usual stacking operation is the special case $\mathcal{F} = \mathcal{N} = \emptyset$.

2.2. Properties of the Construction. Now we want to examine some cases in which the set $\mathcal{R}_{\mathcal{F}, \mathcal{N}}^S(P)$ is non-empty and thus pseudo-stacking can be applied. To simplify the statements of the propositions we introduce some more notation.

DEFINITION 2.5. Let P be a d -polytope and F a facet of P . A subset \mathcal{F} of $\text{adj}(F)$ is called *nonsimple* if there is no pair of adjacent facets $G, G' \in \mathcal{F}$ that have a common $(d-3)$ -face with F .

Nonsimplicity of the set \mathcal{F} implies that there is no simple $(d-3)$ -face (i.e. a $(d-3)$ -face contained in precisely 3 facets) which is a subface of S and the other two facets are in \mathcal{F} .

DEFINITION 2.6. Let P be a d -polytope. A facet F of P is in *bounded position* if the hyperplanes in any subset of $\mathcal{H}_{\text{adj}(F)}$ of cardinality d intersect in a point in H_F^- . Equivalently, F is in bounded position if the intersection of the half-spaces defining P remains bounded if we remove the half-space defined by F .

See Figure 2 for an illustration of the applicability of our definition. The set $\mathcal{F} := \{F_1, F_2\}$ of facets of the polytope P with the chosen simplex facet S in bounded position is nonsimple, while the set $\mathcal{F}' := \{F_2, F_3\}$ is not.

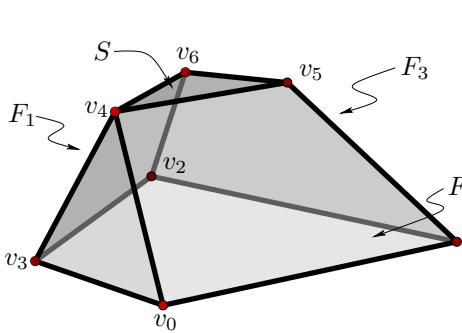


FIGURE 2. A polytope P with simplex facet S in bounded position and $\text{adj}(S) = \{F_1, F_2, F_3\}$.

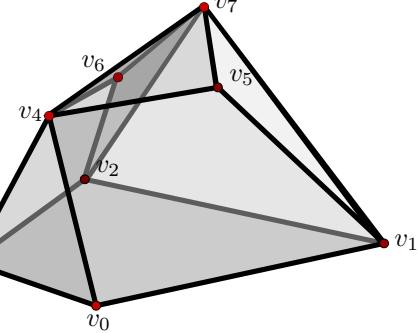


FIGURE 3. The polytope $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ with $\mathcal{F} := \emptyset$ and $\mathcal{N} := \{F_3\}$.

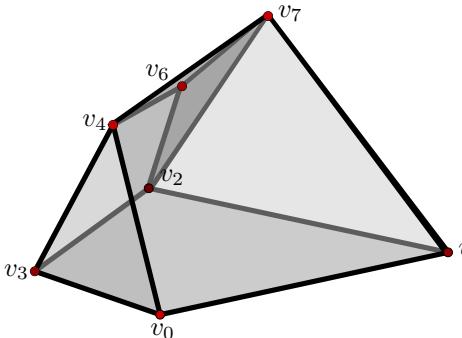


FIGURE 4. The polytope $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ with $\mathcal{F} := \{F_2\}$, $\mathcal{N} := \{F_3\}$.

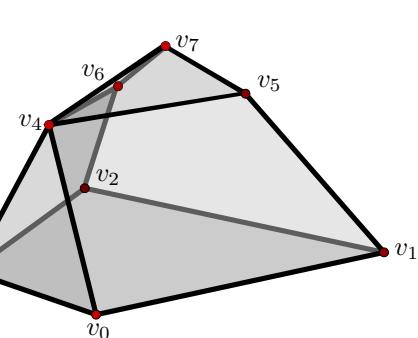


FIGURE 5. The polytope $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ with $\mathcal{F} := \{F_3\}$ and $\mathcal{N} := \emptyset$.

REMARK 2.7. There is always a projective transformation that puts a simplex facet S of a d -polytope P into bounded position, as long as there exists a facet of P not adjacent to S , that is, if P is not a simplex. Without loss of generality we can therefore assume that a simplex facet is in bounded position, if the given polytope itself is not a simplex.

In the following lemma we show that we can always apply our construction to a simplex facet in bounded position, regardless of the choice of \mathcal{F} and \mathcal{N} .

LEMMA 2.8. *Let P be a d -polytope and S a simplex facet of P in bounded position. Choose two disjoint sets $\mathcal{F}, \mathcal{N} \subseteq \text{adj}(S)$. Then the set $\mathcal{R}_{\mathcal{F}, \mathcal{N}}^S(P)$ is not empty.*

PROOF. Let F_1, \dots, F_d be the facets in $\text{adj}(S)$ and number the vertices p_1, \dots, p_d of S in such a way that p_j is the unique vertex not lying in F_j for $j = 1, \dots, d$.

S is a $(d-1)$ -simplex and is in bounded position with respect to P , so all hyperplanes in the set $\mathcal{H}_{\text{adj}(F)}$ intersect in a unique point $p \in H_S^-$. Let $v_j := p_j - p$ for $j = 1, \dots, d$. Define $p' \in \mathbb{R}^d$ by

$$p' := p + \varepsilon \left(\sum_{j: F_j \in \text{adj}(S) \setminus (\mathcal{F} \cup \mathcal{N})} v_j - \sum_{j: F_j \in \mathcal{N}} v_j \right).$$

For $\varepsilon > 0$ small enough p' is contained in $\mathcal{R}_{\mathcal{F}, \mathcal{N}}^S(P)$. \square

The following proposition tells which types of facets occur in $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$.

PROPOSITION 2.9. *Let P be a d -polytope, $d \geq 3$, S a simplex facet of P in bounded position and $\mathcal{F}, \mathcal{N} \subseteq \text{adj}(S)$, $\mathcal{F} \cap \mathcal{N} = \emptyset$. Assume that \mathcal{F} is nonsimple. Then the following properties of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ hold:*

- (1) *The relative interiors of S and of the facets in \mathcal{N} lie in the interior of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$.*
- (2) *Let v be the new vertex added by the pseudo-stacking. $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ has facets of the following four types:*
 - (a) *Facets not in $\{S\} \cup \mathcal{N} \cup \mathcal{F}$ remain unchanged in $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$.*
 - (b) *For any ridge R between S and a facet in $\text{adj}(S) \setminus (\mathcal{F} \cup \mathcal{N})$ we obtain a new facet that is a pyramid over R with apex v .*
 - (c) *For any ridge R between a facet in \mathcal{N} and one in $\text{Fac}(P) \setminus (\mathcal{F} \cup \mathcal{N} \cup \{S\})$ we obtain a new facet that is a pyramid over R with apex v .*
 - (d) *For any facet $F \in \mathcal{F}$ we obtain a facet of the form $\mathbf{PS}_{\emptyset, \mathcal{N}_F}^{S \cap H_F}(F)$ in the pseudo-stacking, where we view F as a polytope in H_F and define $\mathcal{N}_F := \{N \cap F \mid N \in \mathcal{N}\}$.*

PROOF.

- (1) v lies beyond S and all facets in \mathcal{N} by definition. Hence, any ray from v to a point v' in the relative interior of one of these facets intersects P in a segment with boundary points v' and some other point v'' . The segment from v to v'' is contained in $\text{conv}(P, v)$. Hence, v' must be in the interior of $\text{conv}(P, v)$. This proves the first claim.
- (2) The vertices of $\text{conv}(P, v)$ are a subset of the vertices of P and the vertex v . We show that all four facet types mentioned may occur in the pseudo-stacking, and no others.

- (a) The new vertex v lies beneath all facets not in $\mathcal{F} \cup \mathcal{N} \cup \{S\}$. Hence the facet hyperplanes of such facets remain valid and facet defining for the pseudo-stacking of P . This completely describes all facets of $\text{conv}(P, v)$ that do not involve the vertex v .
- (b) Let $F \notin \mathcal{F} \cup \mathcal{N}$ be a facet adjacent to S and $R = F \cap S$. Then v lies beneath H_F and beyond H_S . Hence, any segment between v and a point $v' \in R$ intersects P only in v' ; on the other hand, a segment between v and some $v'' \in F \setminus R$ must intersect the interior of P . In the convex hull of v and P we thus obtain a pyramid over R .
- (c) As above, there is a segment from v to any vertex of a ridge R between a facet $N \in \mathcal{N}$ and a facet $F \in \text{Fac}(P) \setminus (\mathcal{F} \cup \mathcal{N} \cup \{S\})$, but not to any relative interior point of F . Hence, in the convex hull of v and P we again obtain a pyramid over R .
- (d) Let F be a facet in \mathcal{F} and G the ridge it shares with S . We look at F as a polytope defined in the hyperplane H_F . The added vertex v also lies in that hyperplane. G is a facet of F . As \mathcal{F} is nonsimple and S a simplex facet, none of the facets of F adjacent to G is defined by a hyperplane coming from a facet in \mathcal{F} . Let \mathcal{N}_F be the set of ridges F shares with some facet of \mathcal{N} . This is a set of facets of F . Hence, seen as a point in H_F , v is beyond G , beyond all facets in \mathcal{N}_F and beneath all other facet hyperplanes of F . So F is pseudo-stacked by v above G and \mathcal{N}_F .

This completely describes all facets of $\text{conv}(P, v)$ that involve the vertex v . Hence we have described all possible facet types of $\text{conv}(P, v)$. \square

REMARK 2.10. Any simplex facet of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ of type (b) or (c) in Proposition 2.9 is again in bounded position.

The following theorem tells which k -faces of P are also k -faces of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$, for $0 \leq k \leq d - 2$ (cf. also Grünbaum [Grü03, Sec. 5.2]).

THEOREM 2.11. *Let P be a d -polytope, $d \geq 3$, S a simplex facet of P in bounded position and $\mathcal{F}, \mathcal{N} \subseteq \text{adj}(S)$, $\mathcal{F} \cap \mathcal{N} = \emptyset$. Assume that \mathcal{F} is nonsimple.*

A k -face G of P , $0 \leq k \leq d - 2$ is again a k -face of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ if and only if there is a facet $F \in \text{Fac}(P) \setminus (\mathcal{F} \cup \mathcal{N} \cup \{S\})$ that contains G .

PROOF. If G is contained in a facet $F \in \text{Fac}(P) \setminus (\mathcal{F} \cup \mathcal{N} \cup \{S\})$, then F is also a facet of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ by Proposition 2.9, and G is also a face of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$.

So assume that G is a face of P all of whose incident facets are contained in $\mathcal{A} := \mathcal{F} \cup \mathcal{N} \cup \{S\}$. Then any ridge of P containing G is also only incident to facets in \mathcal{A} and we can assume that G is a ridge of P . Because S is a simplex, the intersection of G and S contains a subridge. Hence, by nonsimplicity, there is at most one facet from \mathcal{F} that contains G .

- (1) If there is no facet from \mathcal{F} containing to G , then v lies beyond both facet hyperplanes defining G . Hence, G is in the interior of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$.
- (2) If one of the facets containing G is $F \in \mathcal{F}$ and the other is S , then F is stacked above G , hence G vanishes in the interior of F .
- (3) If one of the facets containing G is $F \in \mathcal{F}$ and the other is in \mathcal{N} , then the pseudo-stacking transforms F into the facet $\mathbf{PS}_{\emptyset, \mathcal{N}'}^{S \cap F}(F)$, where $\mathcal{N}' := \{N \cap F \mid N \in \mathcal{N}\}$ and G is one of the facets in \mathcal{N}' which vanish in the pseudo-stacking. \square

Restricting to the vertices of P we have the following consequence.

COROLLARY 2.12. *Let P be a d -polytope, $d \geq 3$, S a simplex facet of P in bounded position and $\mathcal{F}, \mathcal{N} \subseteq \text{adj}(S)$, $\mathcal{F} \cap \mathcal{N} = \emptyset$. Assume that \mathcal{F} is nonsimple and that any vertex of P is contained in at least one facet $F \in \text{Fac}(P) \setminus (\mathcal{F} \cup \mathcal{N} \cup \{S\})$. Then $f_0(\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)) = f_0(P) + 1$.*

PROOF. By the previous theorem all vertices of P are also vertices of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$. The pseudo-stacking operation adds one new vertex. \square

The next theorems deal with the consequences of the pseudo-stacking operation on k -simpliciality and h -simplicity.

THEOREM 2.13. *Let P be a k -simplicial d -polytope for $d \geq 3$ and $1 \leq k \leq d - 2$ with a simplex facet S in bounded position. Let $\mathcal{F}, \mathcal{N} \subseteq \text{adj}(S)$ be disjoint sets. Assume that \mathcal{F} is nonsimple. Then $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ is k -simplicial.*

PROOF. Let v be the added vertex in the pseudo-stacking. We use the characterisation of the facets in Proposition 2.9. The k -faces in facets that stay unchanged clearly stay combinatorially equivalent to a k -simplex. The k -faces contained in facets that are pyramids over ridges of P are either already faces of P or pyramids over $(k - 1)$ -faces of P with apex v . Hence they are simplices.

All remaining facets are obtained by pseudo-stacking a facet of P . However, only the case $\mathcal{F} = \emptyset$ occurs, that is, in the pseudo-stacking of a facet none of its facets gets stacked. Therefore, the added facets are all pyramids over ridges, which preserve k -simpliciality for $k \leq d - 2$. \square

Now we look more closely at the types of edges that can occur in $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ and determine their number. This is used to establish 2-simplicity for the family of polytopes we construct in the next section. The two cases $\mathcal{N} = \emptyset$ and $\mathcal{N} \neq \emptyset$ are treated separately in the following two propositions, and the latter case is further restricted to $|\mathcal{N}| = 1$.

DEFINITION 2.14. Let P be a polytope. For any face F of P let $\text{fdeg}_P(F)$ be the number of facets of P that contain F .

PROPOSITION 2.15. *Let P be a d -polytope, $d \geq 3$, S a simplex facet of P in bounded position, $\mathcal{F} \subseteq \text{adj}(S)$ nonsimple and $v \in \mathcal{R}_{\mathcal{F}, \emptyset}^S(P)$.*

- (1) *All edges of $\mathbf{PS}_{\mathcal{F}}^S(P)$ are of one of the following two types:*
 - (a) *edges e in P such that $e \subset F$ for a facet $F \in \text{Fac}(P) \setminus (\{S\} \cup \mathcal{F})$;*
 - (b) *edges $e = [v, v']$ for every vertex $v' \in S$; in this case*

$$\text{fdeg}_{\mathbf{PS}_{\mathcal{F}}^S(P)}(e) = \text{fdeg}_S(v').$$

- (2) *$f_1(\mathbf{PS}_{\mathcal{F}}^S(P)) = f_1(P) + d - f$, where $f = |\mathcal{F}|$ if $d = 3$ and $f = 0$ otherwise.*

PROOF.

- (1) The edges of $\mathbf{PS}_{\mathcal{F}}^S(P)$ that are also edges of P are exactly the ones described in (a), by Theorem 2.11. On the other hand, if an edge e of $\mathbf{PS}_{\mathcal{F}}^S(P)$ is not an edge of P , then one of its vertices is v and the other vertex v' is contained in S , which is the only facet of P whose vertices can be seen by v . In this case, every facet of $\mathbf{PS}_{\mathcal{F}}^S(P)$ that contains e corresponds to a ridge $R \subset S$ such that $v' \in R$ — either the facet is of type (b) or of type (d) in Proposition 2.9.

(2) Since S is a simplex there are d edges of type (b). If $d > 3$ then, since \mathcal{F} is nonsimple, every edge of P is contained in at least one facet not in $\mathcal{F} \cup \{S\}$, hence there are $f_1(P)$ edges of type (a). If $d = 3$, then edges are ridges and every ridge $F \cap S$ for $F \in \mathcal{F}$ disappears – namely, it is the base facet for the pseudo-stacking of F (cf. Proposition 2.9). In this case we have $f_1(P) - |\mathcal{F}|$ edges of type (a). \square

PROPOSITION 2.16. *Let P be a d -polytope, $d \geq 3$, S a simplex facet of P in bounded position, $\mathcal{F} \subseteq \text{adj}(S)$ nonsimple, $\mathcal{N} = \{N\}$ with a facet $N \in \text{adj}(S) \setminus \mathcal{F}$ and $v \in \mathcal{R}_{\mathcal{F}, \mathcal{N}}^S(P)$. Additionally, suppose that every vertex of S and N is contained in at least one facet from $\text{Fac}(P) \setminus (\{S\} \cup \mathcal{F} \cup \mathcal{N})$.*

(1) *All edges of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ are of one of the following types:*

- edges e in P such that $e \subset F$ for a facet $F \in \text{Fac}(P) \setminus (\{S\} \cup \mathcal{F} \cup \mathcal{N})$;*
- edges $e = [v, v']$ for every vertex v' of either S or N ; in this case*

$$\text{fdeg}_{\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)}(e) = \text{fdeg}_X(v'),$$

where $X = S$ or $X = N$ respectively;

- edges $e = [v, v']$ for every vertex v' of $S \cap N$.*

(2) $f_1(\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)) = f_1(P) + d + f_0(N) - f_0(S \cap N) - f$, where $f = |\mathcal{F}|$ if $d = 3$ and $f = 0$ otherwise.

PROOF.

- The description of types of edges is complete by the same argument as in the proof before; the only difference is that here the vertex v sees both facets S and N . Also, if v' is a vertex in either S or N , the edge $[v, v']$ is contained in the facets corresponding to the ridges of S , resp. N , that contain v' .
- Again, since S is a simplex, $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ has $d + f_0(N) - f_0(S \cap N)$ edges of type (b) or (c). If $d > 3$ then \mathcal{F} being nonsimple ensures that every edge of P is also an edge of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ of type (a). If $d = 3$, again all edges between facets in \mathcal{F} and S or N disappear; there is no facet in \mathcal{F} that shares ridges with both S and N , since if there were, we had a vertex contained only in facets from $\{S\} \cup \mathcal{F} \cup \mathcal{N}$. Therefore exactly f edges from P are no edges in $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$. \square

We want to show in the next section that certain combinations of pseudo-stacking operations preserve 2-simplicity of a 4-polytope P . For this we have to count the number of facets a subridge of P is in. We do this with the following two propositions. Again, we consider the two cases $\mathcal{N} = \emptyset$ and $\mathcal{N} \neq \emptyset$ separately and restrict the latter to $|\mathcal{N}| = 1$.

PROPOSITION 2.17. *Let P be a d -polytope, $d \geq 3$, S a simplex facet of P in bounded position, $\mathcal{F} \subseteq \text{adj}(S)$ nonsimple and G a subridge of P with $G \subset S$. Define $\varphi := |\{F \in \mathcal{F} \mid G \subset F\}|$.*

Then $0 \leq \varphi \leq 2$ and G is a face of $\mathbf{PS}_{\mathcal{F}}^S(P)$; furthermore,

$$\text{fdeg}_{\mathbf{PS}_{\mathcal{F}}^S(P)}(G) = \text{fdeg}_P(G) + 1 - \varphi.$$

PROOF. G is a ridge of S , hence there are two facets G_1, G_2 of S such that $G_1 \cap G_2 = G$. Since G_1, G_2 are ridges of P , if G is contained in some facet $F \in \mathcal{F}$, then $F \cap S \in \{G_1, G_2\}$. Therefore at most 2 facets in \mathcal{F} can possibly contain G .

If G was not a face of $\mathbf{PS}_{\mathcal{F}}^S(P)$ then by Theorem 2.11 it would be contained in two facets $F_1, F_2 \in \mathcal{F}$. Since S is a simplex, $F_1 \cap F_2$ was a ridge containing G , in contradiction to \mathcal{F} being nonsimple.

Now count the number of facets of $\mathbf{PS}_{\mathcal{F}}^S(P)$ containing G . All facets of P not in $\mathcal{F} \cup \{S\}$ remain facets of $\mathbf{PS}_{\mathcal{F}}^S(P)$ by Proposition 2.9; therefore, $\text{fdeg}_P(G) - \varphi - 1$ facets still contain G in $\mathbf{PS}_{\mathcal{F}}^S(P)$. All facets in \mathcal{F} containing G also stay facets of $\mathbf{PS}_{\mathcal{F}}^S(P)$ containing G . Additionally, for every one of the two ridges R of P with $G \subset R \subset S$, we either pseudo-stacked the facet F with $F \cap S = R$ (if it is in \mathcal{F}) or we get a new facet, which is a pyramid over R ; in the latter case, we have not counted it yet.

In total, there are $\text{fdeg}_P(G) - \varphi - 1 + \varphi + 2 - \varphi = \text{fdeg}_P(G) + 1 - \varphi$ facets of $\mathbf{PS}_{\mathcal{F}}^S(P)$ that contain G . \square

PROPOSITION 2.18. *Let P , S , \mathcal{F} , G and φ as in Proposition 2.17. Choose a facet $N \in \text{adj}(S) \setminus \mathcal{F}$ and set $\mathcal{N} = \{N\}$.*

(1) *If $G \subset S$ and $G \not\subset N$, then G is a face of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ and*

$$\text{fdeg}_{\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)}(G) = \text{fdeg}_P(G) + 1 - \varphi.$$

(2) *If $G \subset N$ and every facet containing G is either S , N or in \mathcal{F} , then G is not a face of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$; otherwise, G is a face of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ and*

$$\text{fdeg}_{\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)}(G) = \text{fdeg}_P(G) + 1 - \varepsilon - \varphi,$$

where $\varepsilon = 1$ if $G \subset S$ and $\varepsilon = 0$ if not.

PROOF. (1) can be shown in the same way as Proposition 2.17, since the facet N has no influence whatsoever on G .

Suppose G is contained in N . By Theorem 2.11, G is a face of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ if and only if there is facet containing G that is not in $\mathcal{F} \cup \{S, N\}$. In this case, the number of facets of $\mathbf{PS}_{\mathcal{F}, \mathcal{N}}^S(P)$ can be computed in the same way as in the previous proof, except that if $G \subset S$, an additional facet containing G (namely S) disappears. \square

REMARK 2.19. The construction given in this section can be further generalised:

- (1) Nonsimplicity of the set \mathcal{F} is not necessary.
- (2) The facet S need not be a simplex.
- (3) In this case the construction can be adapted to only apply to a part of S . Several new families of polytopes arise via these generalisations. However, keeping track of the changes in the f -vector is more subtle in these cases. We do not need them here, so we omit their treatment.

3. Generating 2-simple and 2-simplicial 4-polytopes

In this part we use the pseudo-stacking operation defined in the previous section to extend elementary 2-simple and 2-simplicial 4-polytopes while maintaining these properties. As we have seen in Theorem 2.13, pseudo-stacking preserves 2-simpliciality. The more difficult part is to also preserve 2-simplicity; we accomplish this by applying the operation five, resp. four times in a suitable way, such that edges whose degree initially increases are again contained in only 3 facets at the end.

To simplify the task of keeping track of the facets generated in each step of the construction we introduce two ways of distinguishing certain facets.

TABLE 2. Summary of parameters and involved faces for the first construction.

step i	base facet S_i	facets in \mathcal{F}_i	new facets	changing edges	degree
1 (Fig. 6(b))	S	$F_S(e_0, e_1, e_2)$	$F(v_0, v_1, v_3, v_4)$ $F(v_1, v_2, v_3, v_4)$ $F(v_0, v_2, v_3, v_4)$ $e_6 := [v_0, v_4]$ $e_7 := [v_1, v_4]$ $e_8 := [v_2, v_4]$ $e_9 := [v_2, v_4]$	$e_6 := [v_0, v_4]$ $e_7 := [v_1, v_4]$ $e_8 := [v_2, v_4]$ $e_9 := [v_2, v_4]$ e_3 e_4 e_5	3 3 3 3 $d_3 + 1$ $d_4 + 1$ $d_5 + 1$
2 (Fig. 6(c))	$F(v_0, v_1, v_3, v_4)$	$F(v_1, v_2, v_3, v_4)$ $F_{S_2}(e_0, e_3, e_4)$	$F(v_0, v_3, v_4, v_5)$ $F(v_0, v_1, v_4, v_5)$ $[v_0, v_5]$ $[v_1, v_5]$ $[v_3, v_5]$ $[v_4, v_5]$ e_4 e_6	$[v_0, v_5]$ $[v_1, v_5]$ $[v_3, v_5]$ $[v_4, v_5]$ d_4 d_5	3 3 3 3 4
3 (Fig. 6(d))	$F(v_0, v_2, v_3, v_4)$	$F(v_0, v_3, v_4, v_5)$ $F_{S_3}(e_2, e_3, e_5)$	$F(v_2, v_3, v_4, v_6)$ $F(v_0, v_2, v_4, v_6)$ $e_{14} := [v_0, v_6]$ $e_{15} := [v_2, v_6]$ $e_{16} := [v_3, v_6]$ $e_{17} := [v_4, v_6]$ e_3 e_8	$e_{14} := [v_0, v_6]$ $e_{15} := [v_2, v_6]$ $e_{16} := [v_3, v_6]$ $e_{17} := [v_4, v_6]$ d_3 d_5	3 3 3 3 4
4 (Fig. 6(e))	$F(v_2, v_3, v_4, v_6)$	$F(v_1, v_2, v_3, v_4)$ $F_{S_4}(e_5, e_{16}, e_{17})$	$F(v_3, v_4, v_6, v_7)$ $F(v_2, v_4, v_6, v_7)$ $[v_2, v_7]$ $[v_3, v_7]$ $[v_4, v_7]$ $[v_6, v_7]$ e_5 e_{17}	$[v_2, v_7]$ $[v_3, v_7]$ $[v_4, v_7]$ $[v_6, v_7]$ d_5 d_5	3 3 3 3 4
5 (Fig. 6(f))	$F(v_0, v_2, v_4, v_6)$	$F(v_0, v_3, v_4, v_5, v_6)$ $F(v_2, v_4, v_6, v_7)$ $F_{S_5}(e_2, e_6, e_8)$	$F(v_0, v_2, v_6, v_8)$ $[v_0, v_8]$ $[v_2, v_8]$ $[v_4, v_8]$ $[v_6, v_8]$ e_6 e_8 e_{17}	$[v_0, v_8]$ $[v_2, v_8]$ $[v_4, v_8]$ $[v_6, v_8]$ d_3 d_3 d_5	3 3 3 3 3

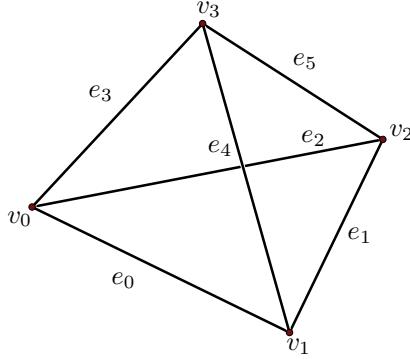
DEFINITION 3.1. Let P be a d -polytope and S a facet of P .

- (1) If the vertices v_0, \dots, v_k all lie in one facet of P and define this facet uniquely, then we denote this facet by $F(v_0, \dots, v_k)$.
- (2) If e_0, \dots, e_m are edges of a facet G of S , then we denote the unique facet of P adjacent to S via G by $F_S(e_0, \dots, e_m)$.

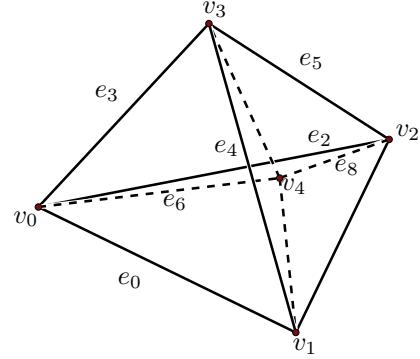
3.1. The first construction. Let P be a 4-polytope and S a simplex facet of P in bounded position. S has six edges and four vertices, which we label e_0, e_1, \dots, e_5 and v_0, v_1, v_2, v_3 according to Figure 6(a).

We proceed with the five pseudo-stacking steps described in Table 2. In Step i we construct the polytope $P^{(i)} := \mathbf{PS}_{\mathcal{F}_i}^{S_i}(P^{(i-1)})$, with $P^{(0)} := P$ and the parameters as given in Table 2. For every step we additionally give a complete list of facets of type (b) in Proposition 2.9, as well as of edges that either show up by Proposition 2.15 or change their degree by Proposition 2.17; note that the only edges of P that change degree during the process are e_3, e_4 and e_5 , whose degrees in P are denoted by d_3, d_4 and d_5 respectively.

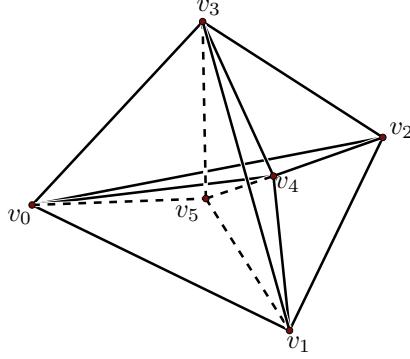
Due to Remark 2.10 the base facets in each step are in bounded position; additionally, it can be verified easily that the sets \mathcal{F}_i are nonsimple.



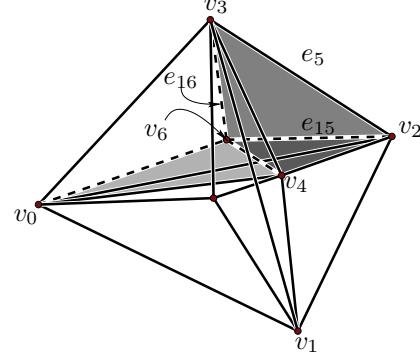
(a) Labelling of the vertices and edges of the simplex S .



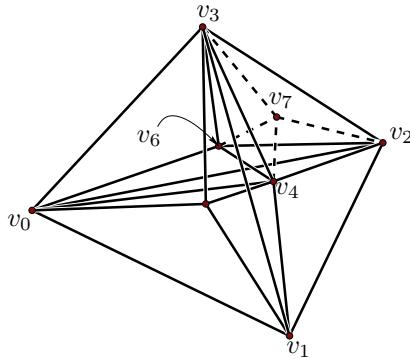
(b) Part of the Schlegel diagram of $P^{(1)}$; the new edges e_6, \dots, e_9 (partly labelled) are drawn as dashed lines.



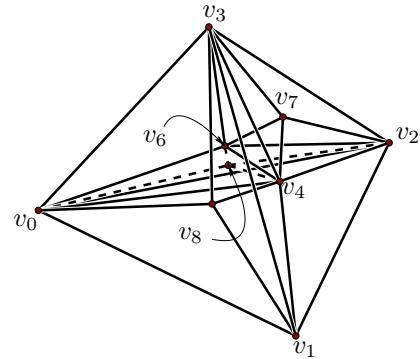
(c) Part of the Schlegel diagram of $P^{(2)}$; note that $v_1 \dots, v_5$ define a bipyramid over a triangle, which occurred by stacking the facet $F(v_1, v_2, v_3, v_4)$ of $P^{(1)}$.



(d) Part of the Schlegel diagram of $P^{(3)}$; the shaded facets are the base facets for the last two pseudo-stacking steps.



(e) Part of the Schlegel diagram of $P^{(4)}$ showing the pseudo-stacking of the facet $F(v_2, v_3, v_4, v_6)$ of $P^{(3)}$.



(f) Part of the Schlegel diagram of $P^{(5)}$ showing the pseudo-stacking of the facet $F(v_0, v_2, v_4, v_6)$ of $P^{(4)}$.

FIGURE 6. The first construction.

DEFINITION 3.2. Let P be a 4-polytope with a simplex facet S in bounded position. We denote the polytope $P^{(5)}$ obtained by applying the five steps of Table 2 by $\mathcal{I}^1(P; S)$.

REMARK 3.3. Note that $\mathcal{I}^1(P; S)$ implicitly depends on a labelling of the vertices of S . Choosing different labellings may result in combinatorially different polytopes for the same choice of P and S .

LEMMA 3.4. *Let P be a 4-polytope with a simplex facet S in bounded position and the vertices of S numbered in arbitrary order. Then $\mathcal{I}^1(P; S)$ has again simplex facets in bounded position. Additionally, all edges of P are still present in $\mathcal{I}^1(P; S)$, their degrees remain unchanged and all new edges have degree 3. \square*

3.2. A second construction. Let Q be a 4-polytope with a simplex facet S in bounded position. We label the vertices and edges of S in the same way as in the previous section (see Figure 6(a)).

The construction is described by Table 3, in the same way as before. The main difference is that in Step 3 the set \mathcal{N}_3 is not empty, so in general we construct $Q^{(i)} := \mathbf{PS}_{\mathcal{F}_i, \mathcal{N}_i}^{S_i}(Q^{(i-1)})$ with the given parameters. Note that one effect of this is that in Step 3 one edge disappears by Theorem 2.11; also, the only edges of Q changing their degree are e_1, e_2 and e_5 , whose degrees are labelled accordingly.

DEFINITION 3.5. Let Q be a 4-polytope with a simplex facet S in bounded position. We denote the polytope $Q^{(4)}$ obtained by applying the four steps of Table 3 by $\mathcal{I}^2(Q; S)$.

Remark 3.3 also applies to $\mathcal{I}^2(P; S)$. As in the first construction, we have again created simplex facets in bounded position in $\mathcal{I}^2(Q; S)$ and 2-simplicity is preserved.

LEMMA 3.6. *Let Q be a 4-polytope with a simplex facet S in bounded position and the vertices of S numbered in arbitrary order. Then the simplex facet of $\mathcal{I}^2(Q; S)$ constructed in the last step is in bounded position. Furthermore, all edges of Q are again edges of $\mathcal{I}^2(Q; S)$ with the same degree, and $f\text{deg}(e) = 3$ for all edges e in $\mathcal{I}^2(Q; S)$ that were not edges in Q . \square*

3.3. Properties of the Constructions. The two constructions preserve the properties we are interested in.

THEOREM 3.7. *Let P be a 4-polytope with a simplex facet S in bounded position. Then*

$$f(\mathcal{I}^1(P; S)) = f(P) + (5, 20, 20, 5),$$

and

$$f(\mathcal{I}^2(P; S)) = f(P) + (4, 16, 16, 4).$$

PROOF. In every step of both constructions, we add one vertex each by Corollary 2.12. Hence $f_0(\mathcal{I}^1(P; S)) = f_0(P) + 5$ and $f_0(\mathcal{I}^2(P; S)) = f_0(P) + 4$.

Propositions 2.15 and 2.16 tell how many edges are added in each step. Accordingly, the first construction adds 4 edges in each step. In Steps 1, 2 and 4 of the second construction, we also add 4 edges each; in the third step 5 edges are added, but one edge is destroyed, so the overall change is also 4. Therefore, we have $f_1(\mathcal{I}^1(P; S)) = f_1(P) + 5 \cdot 4 = f_1(P) + 20$ and $f_1(\mathcal{I}^2(P; S)) = f_1(P) + 4 \cdot 4 = f_1(P) + 16$. Finally, the number of new facets in the respective steps

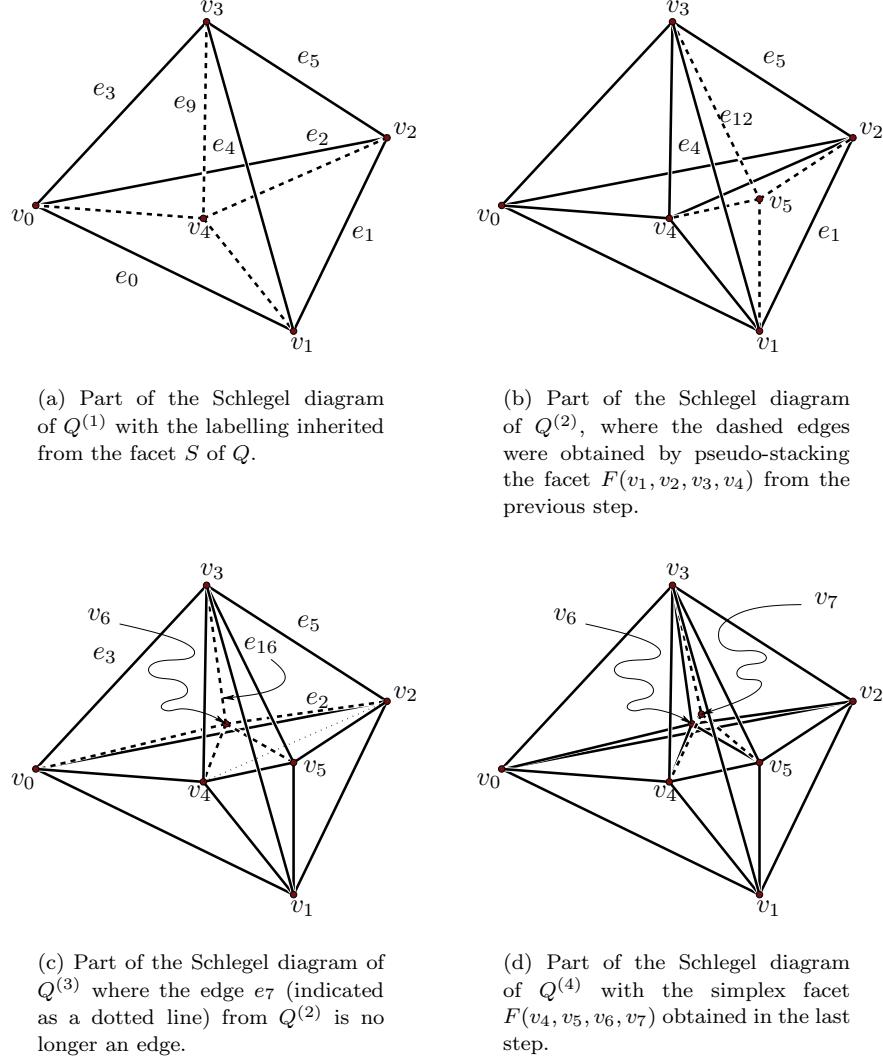


FIGURE 7. The second construction.

can be read off from Tables 2 and 3 (or Proposition 2.9) – note that the facets S_i in every step, as well as the facet of $Q^{(3)}$ in \mathcal{N}_3 , disappear. Summarising, we get $f_3(\mathcal{I}^1(P; S)) = f_3(P) + 5$ and $f_3(\mathcal{I}^2(P; S)) = f_3(P) + 4$.

By Euler's equation, the number of ridges is determined by f_0, f_1 and f_3 , which implies the claim. \square

COROLLARY 3.8. *Let P be a 4-polytope with a simplex facet S in bounded position. Then $g_2(\mathcal{I}^1(P; S)) = g_2(\mathcal{I}^2(P; S)) = g_2(P)$.*

PROOF. Define $a(P) := f_{02}(P) - 3f_2(P)$ and $b(P) := f_1(P) - 4f_0(P)$. Then $g_2(P) = a(P) + b(P) + 10$.

TABLE 3. Parameters describing the second construction.

step i	base facet S_i	facets in \mathcal{F}_i	facets in \mathcal{N}_i	new facets	changing edges	degree
1 (Fig. 7(a))	S	$F_S(e_0, e_3, e_4)$	—	$F(v_0, v_1, v_2, v_4)$ $F(v_1, v_2, v_3, v_4)$ $F(v_0, v_2, v_3, v_4)$ $e_9 := [v_2, v_4]$ e_1 e_2 e_5	$[v_0, v_4]$ $[v_1, v_4]$ $[v_2, v_4]$ $d_1 + 1$ $d_2 + 1$ $d_5 + 1$	3 3 3 3 $d_1 + 1$ $d_2 + 1$ $d_5 + 1$
2 (Fig. 7(b))	$F(v_1, v_2, v_3, v_4)$	$F(v_0, v_1, v_2, v_4)$ $F_{S_2}(e_1, e_4, e_5)$	—	$F(v_1, v_3, v_4, v_5)$ $F(v_2, v_3, v_4, v_5)$ $e_{12} := [v_3, v_5]$ $[v_4, v_5]$ e_1 e_9	$[v_1, v_5]$ $[v_2, v_5]$ $[v_3, v_5]$ $[v_4, v_5]$ d_1 4	3 3 3 3 d_1 4
3 (Fig. 7(c))	$F(v_0, v_2, v_3, v_4)$	$F(v_0, v_1, v_2, v_4, v_5)$ $F_{S_3}(e_2, e_3, e_5)$	$F(v_2, v_3, v_4, v_5)$	$F(v_0, v_3, v_4, v_6)$ $F(v_2, v_3, v_5, v_6)$ $F(v_3, v_4, v_5, v_6)$ $e_{16} := [v_3, v_6]$ $[v_4, v_6]$ $[v_5, v_6]$ e_2 e_5 e_{12}	$[v_0, v_6]$ $[v_2, v_6]$ $[v_3, v_6]$ $[v_4, v_6]$ $[v_5, v_6]$ d_2 d_5 4	3 3 4 3 3 d_2 d_5 4
4 (Fig. 7(d))	$F(v_3, v_4, v_5, v_6)$	$F(v_0, v_3, v_4, v_6)$ $F(v_1, v_3, v_4, v_5)$ $F(v_2, v_3, v_5, v_6)$	—	$F(v_4, v_5, v_6, v_7)$	$[v_3, v_7]$ $[v_4, v_7]$ $[v_5, v_7]$ $[v_6, v_7]$ e_9 e_{12} e_{16}	3 3 3 3 3 3 3

We first show that $a(\mathcal{T}^1(P; S)) = a(P)$. Any 2-face of $\mathcal{T}^1(P; S)$ that is not a face of P is a triangle by Proposition 2.9. Furthermore, any 2-face of P that is not a face of $\mathcal{T}^1(P; S)$ is a ridge of the base facet in one of the five steps; since all base facets are simplices, all such 2-faces are also triangles. Therefore, none of these faces causes a change in $a(P)$.

For $b(P)$, we have, by Theorem 3.7,

$$\begin{aligned} b(\mathcal{I}^1(P; S)) &= f_1(\mathcal{I}^1(P; S)) - 4f_0(\mathcal{I}^1(P; S)) \\ &= f_1(P) + 20 - 4(f_0(P) + 5) = f_1(P) - 4f_0(P) = b(P). \end{aligned}$$

This implies $g_2(\mathcal{I}^1(P; S)) = g_2(P)$.

Similar reasoning shows the claim for $\mathcal{I}^2(P; S)$. \square

THEOREM 3.9. *Let P be a 2-simplicial, 2-simple 4-polytope with a simplex facet S in bounded position. Then $\mathcal{I}^1(P; S)$ and $\mathcal{I}^2(P; S)$ are again 2-simple and 2-simplicial.*

PROOF. By Theorem 2.13, all intermediate polytopes in both constructions are 2-simplicial, hence also $\mathcal{I}^1(P; S)$ and $\mathcal{I}^2(P; S)$. Furthermore, all edges in $\mathcal{I}^1(P; S)$ and $\mathcal{I}^2(P; S)$ have degree 3 by Lemmas 3.4 and 3.6, and since this is true for P . \square

REMARK 3.10. There are more sequences of pseudo-stacking operations that preserve $g_2 = 0$ than the two given here. Even more can be constructed with the help of a further generalised pseudo-stacking operation, see Remark 2.19.

A more symmetric version using these generalisations also applies in dimensions $d \geq 5$, i.e. using this construction one can obtain elementary 2-simple and 2-simplicial d -polytopes with arbitrarily large numbers of vertices.

4. Examples and Results

Now we are ready to prove our main result: that ℓ_1 is in fact an extremal ray of \mathcal{C} . We prove this by providing three examples P_9 , P_{10} , and P_{11} of 2-simple, 2-simplicial elementary 4-polytopes with 9, 10, and 11 vertices that have simplex facets in bounded position. Using these as input for Theorem 3.7 we obtain such polytopes for arbitrarily high numbers of vertices.

4.1. Examples. For each of the following three examples we give their vertex-facet incidences together with a Schlegel diagram. Explicit rational coordinates are in Table 4. Calculations for the realisations were done by computer with the **polymake** system [GJ00]. Clients producing the examples below and applying the constructions, as well as coordinates for many more elementary 2-simple and 2-simplicial 4-polytopes are available from the authors.

The combinatorial descriptions of the examples, together with many more such polytopes having a small number of vertices, were found via a complete enumeration approach using a client for the **polymake** system.

- (1) The smallest non-trivial 2-simplicial, 2-simple 4-polytope is P_9 (see also the remark below). See Figure 8.
- (2) The second example, P_{10} , has 10 vertices. See Figure 9. Note that P_{10} is not combinatorially equivalent to the 4-dimensional hypersimplex, nor its dual – although it has the same flag vector. Nevertheless, the hypersimplex itself can be obtained with the extended construction described in Remark 2.19.
- (3) The last example, P_{11} , has 11 vertices. See Figure 10.

REMARK 4.1. P_9 and P_{10} are both self-dual. P_9 is the unique non-trivial 2-simple and 2-simplicial 4-polytope with smallest number of vertices, i.e. except for the simplex Δ_4 there are no further 2-simplicial, 2-simple 4-polytopes with less than 10 vertices. We do not prove this here; part of the proof can be found in [Paf05].

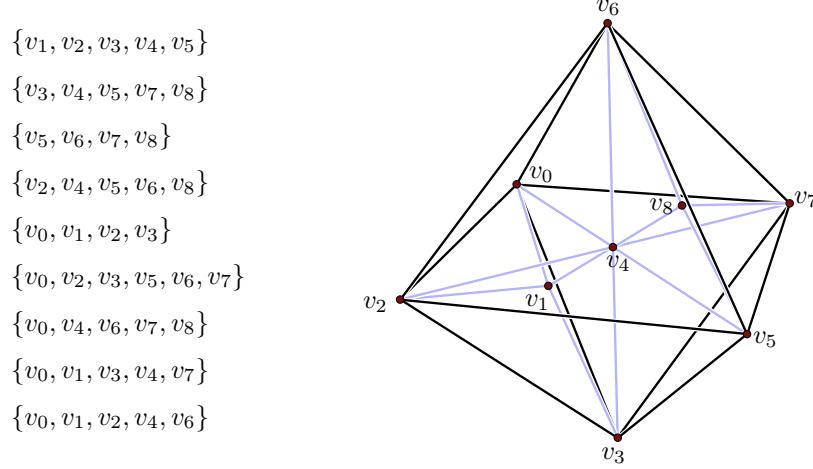


FIGURE 8. The vertex-facet-incidences and a Schlegel diagram of the polytope P_9 .

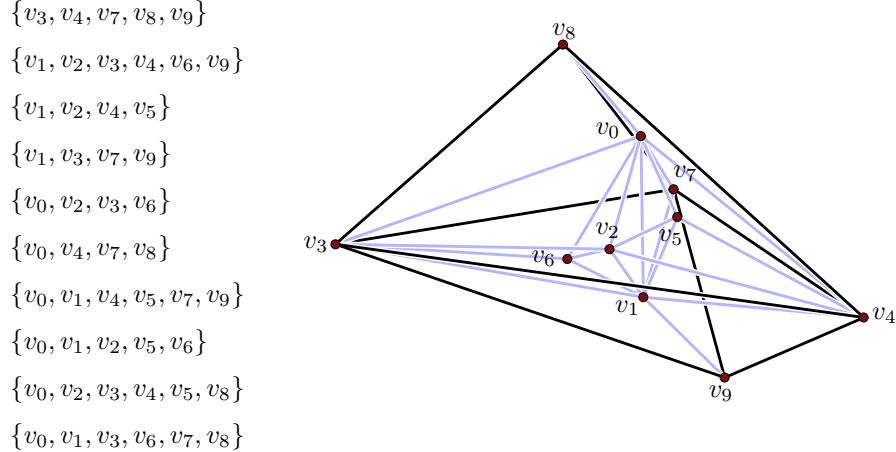


FIGURE 9. The vertex-facet-incidences and a Schlegel diagram of the polytope P_{10} .

REMARK 4.2. The two examples P_9 and P_{10} can in fact themselves be obtained by applying the two constructions to the 4-simplex: $P_9 = \mathcal{I}^2(\Delta_4; S)$ and $P_{10} = \mathcal{I}^1(\Delta_4; S)$ with an arbitrary facet S of Δ_4 . Note however, that the simplex has no facet in bounded position; nevertheless, the constructions remain valid, since in both cases the first step requires the added vertex to lie in only one adjacent facet hyperplane.

Also, P_{11} can be obtained via three pseudo-stacking steps in the following way. Take a 3-dimensional octahedron O^3 . Choose a facet R of O^3 and let R_1, R_2 , and R_3 be the facets adjacent to R . Let B^3 be the polytope obtained from O^3 by stacking the facet R . Let PB^3 be the pyramid over B^3 with apex v and F_1, F_2 ,

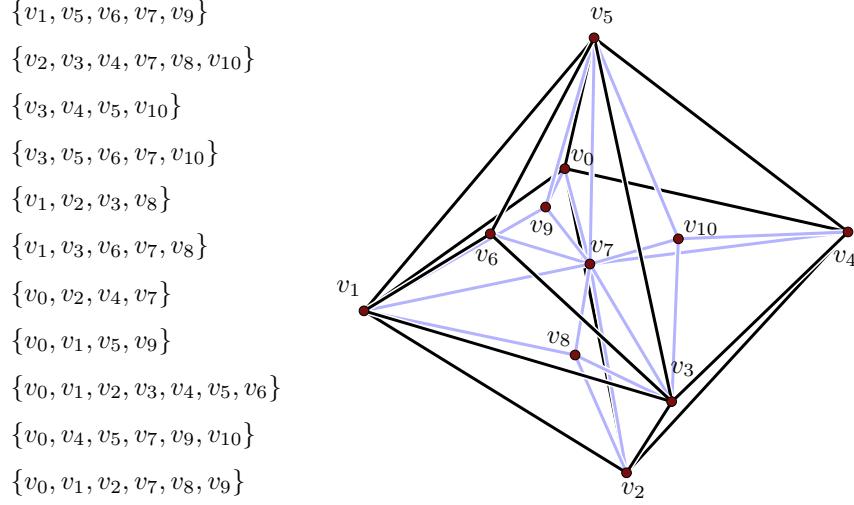


FIGURE 10. The vertex-facet-incidences and a Schlegel diagram of the polytope P_{11} .

and F_3 be the facets of PB^3 arising as pyramids over R_1 , R_2 , and R_3 . PB^3 is an elementary 2-simplicial polytope, but it is not 2-simple.

Let \mathcal{F}_i be the sets of facets adjacent to F_i containing the vertex v . Then $|\mathcal{F}_i| = 3$. Set $P^{(0)} := PB^3$ and define polytopes $P^{(i)} := \mathbf{PS}_{\mathcal{F}_i}^{F_i}(P^{(i-1)})$ for $i = 1, 2, 3$. Then $P_{11} = P^{(3)}$. Note that some facets in \mathcal{F}_2 and \mathcal{F}_3 are stacked by the previous steps, but this does not influence the result.

P_{11} demonstrates that one can obtain elementary 2-simple and 2-simplicial 4-polytopes by pseudo-stacking polytopes without these properties in a suitable way.

4.2. Conclusions. With these examples and the constructions of Section 3, 2-simplicial and 2-simple 4-polytopes with $g_2 = 0$ can be found for almost all numbers of vertices. Here is the main theorem of this paper.

THEOREM 4.3. *Elementary 2-simplicial, 2-simple 4-polytopes with k vertices exist for $k = 5, 9, 10, 11$ and $k \geq 13$.*

PROOF. We show that for k as given in the claim there exist elementary 2-simplicial, 2-simple 4-polytopes with k vertices that have at least one simplex facet in bounded position.

For $k = 9, 10, 11$, the above examples P_9 , P_{10} , resp. P_{11} have the desired properties. Let $k \geq 13$. By induction there is such a polytope P with $k - 5$ or $k - 4$ vertices and a simplex facet S in bounded position. Then the polytope $\mathcal{I}^1(P; S)$, resp. $\mathcal{I}^2(P; S)$ are elementary 2-simplicial, 2-simple 4-polytopes on k vertices by Lemma 3.4 resp. 3.6 and Corollary 3.8. \square

COROLLARY 4.4. *The ray ℓ_1 is contained in the convex hull of all flag vectors of 4-polytopes.* \square

There are also many 2-simplicial, 2-simple 4-polytopes with $g_2 > 0$; still the existence of a 2-simplicial, 2-simple 4-polytope with 12 vertices is an open question.

TABLE 4. Coordinates for the three examples of 2-simple and 2-simplicial 4-polytopes that we have discussed in this section.

P_9	P_{10}	P_{11}
(3, 0, 0, 0)	(-9, -3, -3, -3)	(1, 0, 0, 0)
(1, 1, 1, 1)	(-3, 9, -3, -3)	(0, 1, 0, 0)
(0, 3, 0, 0)	(-3, -3, -3, -3)	(0, 0, 1, 0)
(0, 0, 3, 0)	(-3, -3, 9, -3)	(-1, 0, 0, 0)
(0, 0, 0, $\frac{3}{2}$)	(-3, -3, -3, 9)	(0, -1, 0, 0)
(-3, 0, 0, 0)	(1, -3, -7, 1)	(0, 0, -1, 0)
(0, 0, -3, 0)	(-3, 1, 1, -7)	(- $\frac{11}{21}$, $\frac{11}{21}$, - $\frac{11}{21}$, 0)
(0, -3, 0, 0)	(3, 3, 3, 3)	(- $\frac{11}{147}$, $\frac{11}{147}$, - $\frac{11}{147}$, 1)
(-1, -1, -1, 1)	(5, -3, 5, 1)	(- $\frac{428}{1617}$, $\frac{428}{1617}$, $\frac{68}{147}$, $\frac{1}{2}$)
	(-3, 5, 1, 5)	($\frac{68}{147}$, $\frac{428}{1617}$, - $\frac{428}{1617}$, $\frac{1}{2}$)
		(- $\frac{428}{1617}$, - $\frac{68}{147}$, - $\frac{428}{1617}$, $\frac{1}{2}$)

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